

# **ns-3 Introduction**

**Tom Henderson (University of Washington)**

**July 2014**

# Agenda

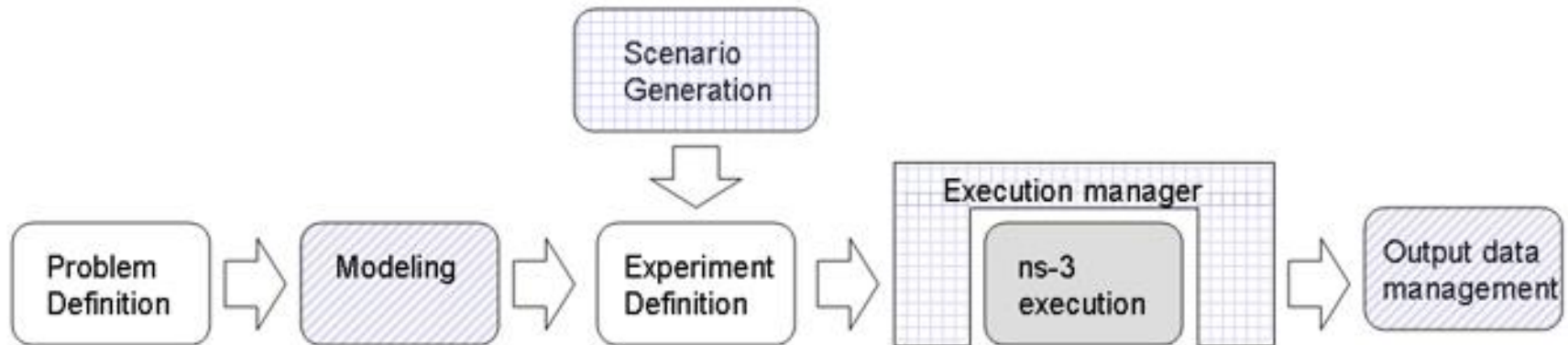
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- ns-3 project overview
  - What is ns-3?
  - Why use ns-3?
  - Project organization
  - Relationship to ns-2
  - Future directions
- Getting started with ns-3

# Discrete event network simulator

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- Model of the evolution of a networked system through discrete events in time
- Used for experimentation and education



# ns-3 simulation basics

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- Simulation time advances in discrete jumps from event to event
- C++ functions schedule events to occur at specific simulation times
- A simulation scheduler orders the event execution
- `Simulation::Run()` gets it all started
- Simulation stops at specific time or when events end

# Software overview

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- ns-3 is written in C++, with bindings available for Python
  - simulation programs are C++ executables or Python programs
  - ~350,000 lines of C++ (estimate based on cloc source code analysis)
- ns-3 is a GNU GPLv2-licensed project
- ns-3 is mainly supported for Linux, OS X, and FreeBSD
  - Windows Visual Studio port available
- ns-3 is not backwards-compatible with ns-2

# Software orientation

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Key differences from other network simulators:

1) Command-line, Unix orientation

– vs. Integrated Development Environment (IDE)

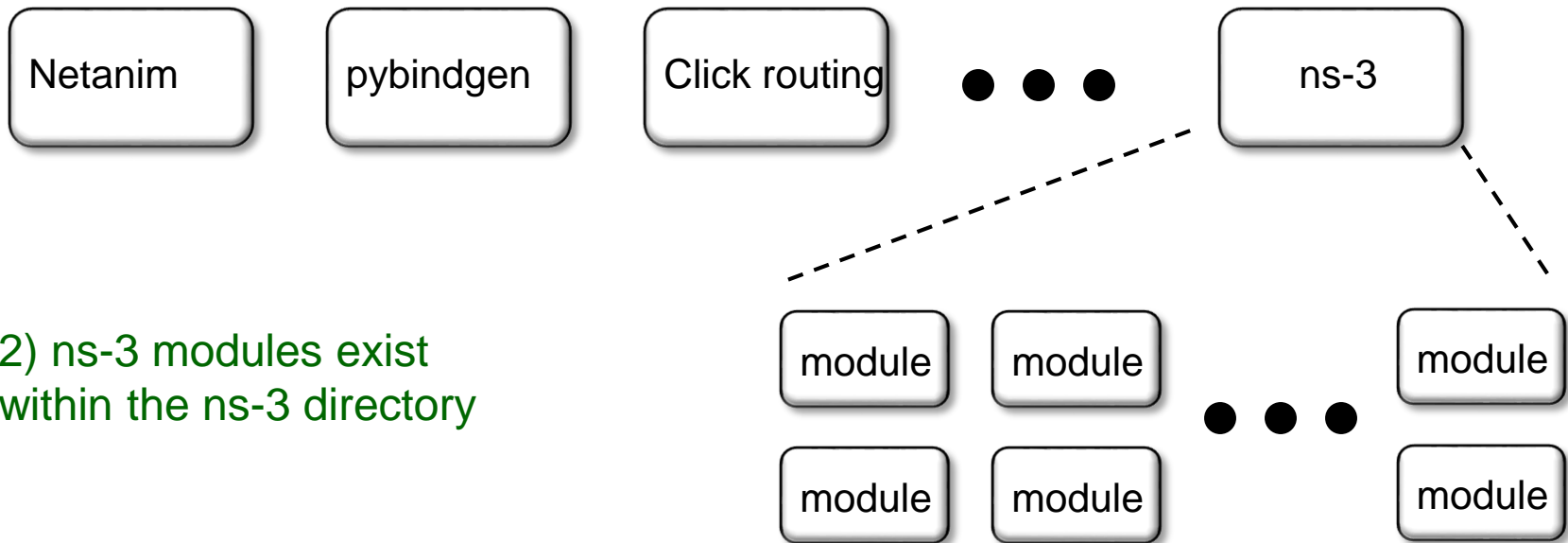
2) Simulations and models written directly in C++ and Python

– vs. a domain-specific simulation language

# Software organization

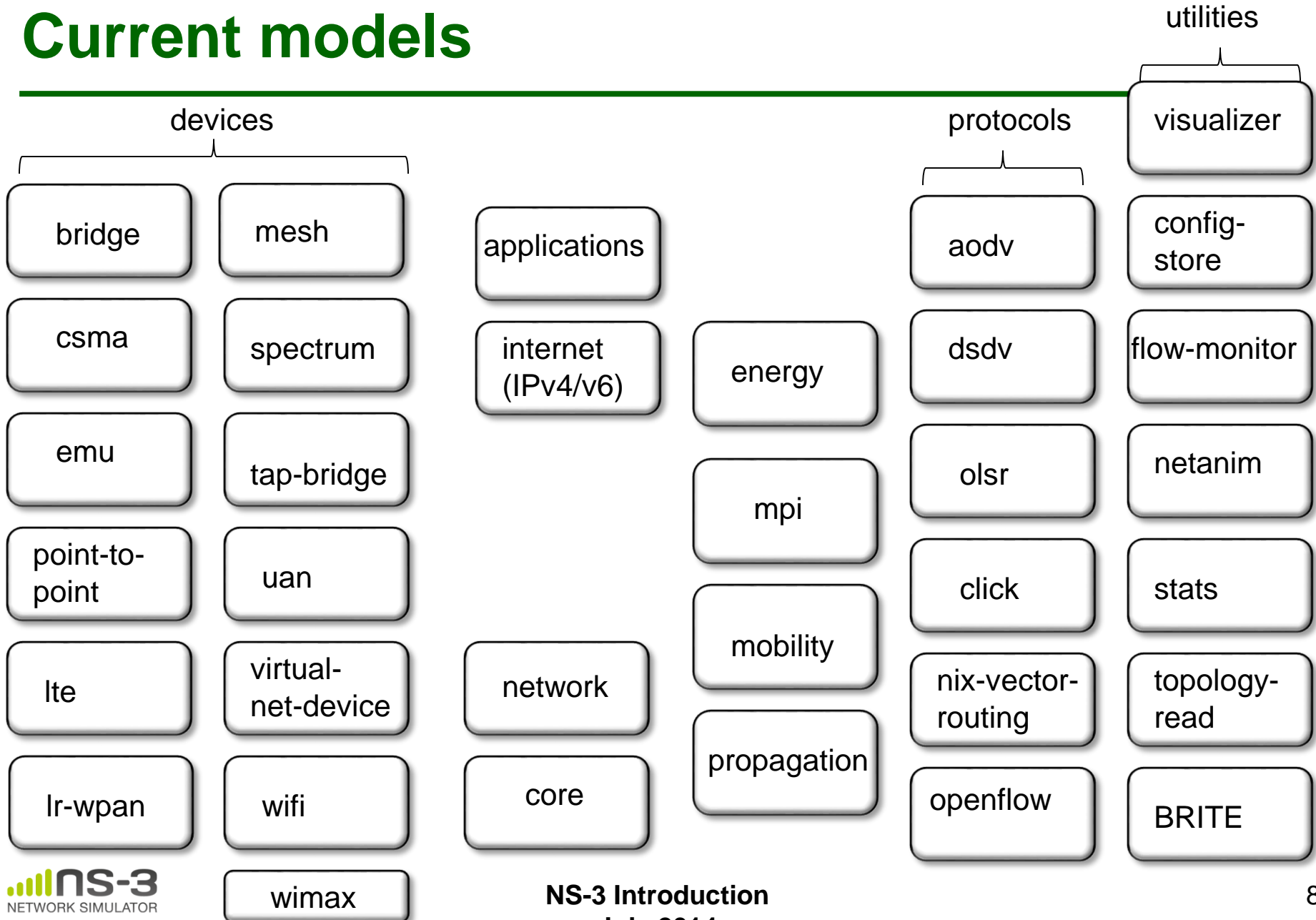
- Two levels of ns-3 software and libraries

1) Several supporting libraries, not system-installed, can be in parallel to ns-3



2) ns-3 modules exist within the ns-3 directory

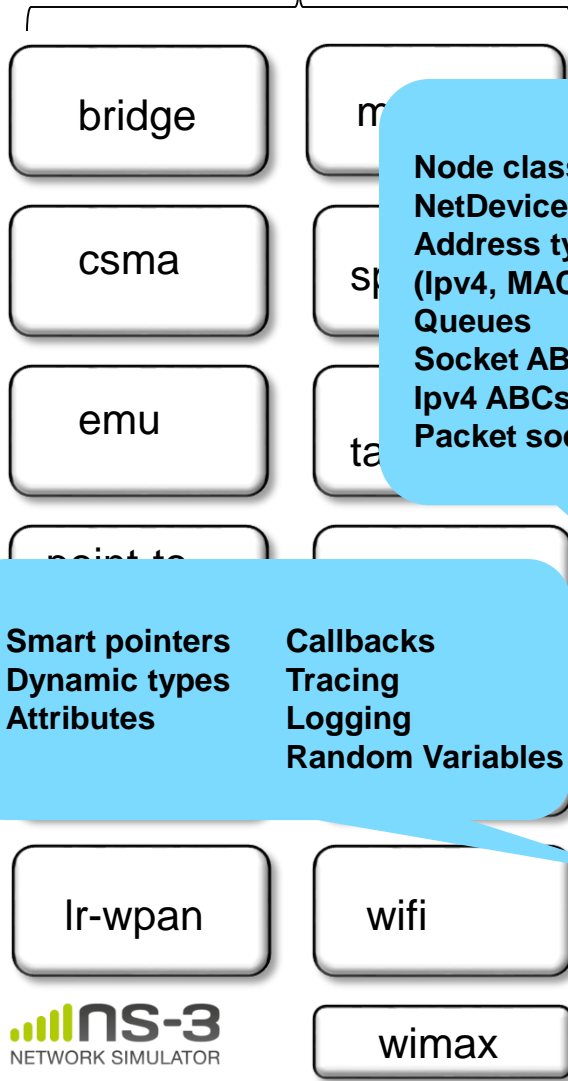
# Current models





# Current models

devices



Node class  
 NetDevice ABC  
 Address types  
 (Ipv4, MAC, etc.)  
 Queues  
 Socket ABC  
 Ipv4 ABCs  
 Packet sockets

Smart pointers  
 Dynamic types  
 Attributes

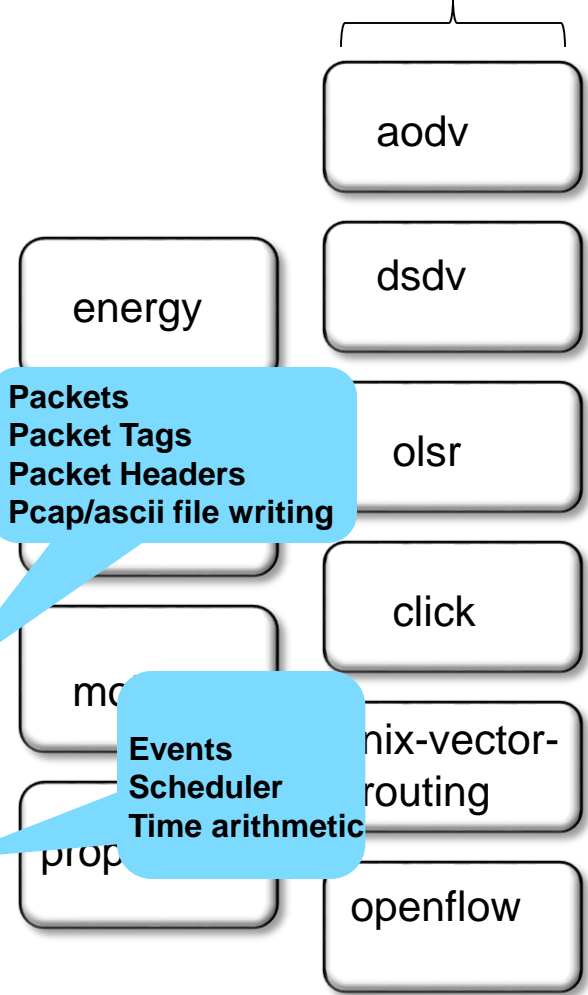
Callbacks  
 Tracing  
 Logging  
 Random Variables



Packets  
 Packet Tags  
 Packet Headers  
 Pcap/ascii file writing

Events  
 Scheduler  
 Time arithmetic

protocols

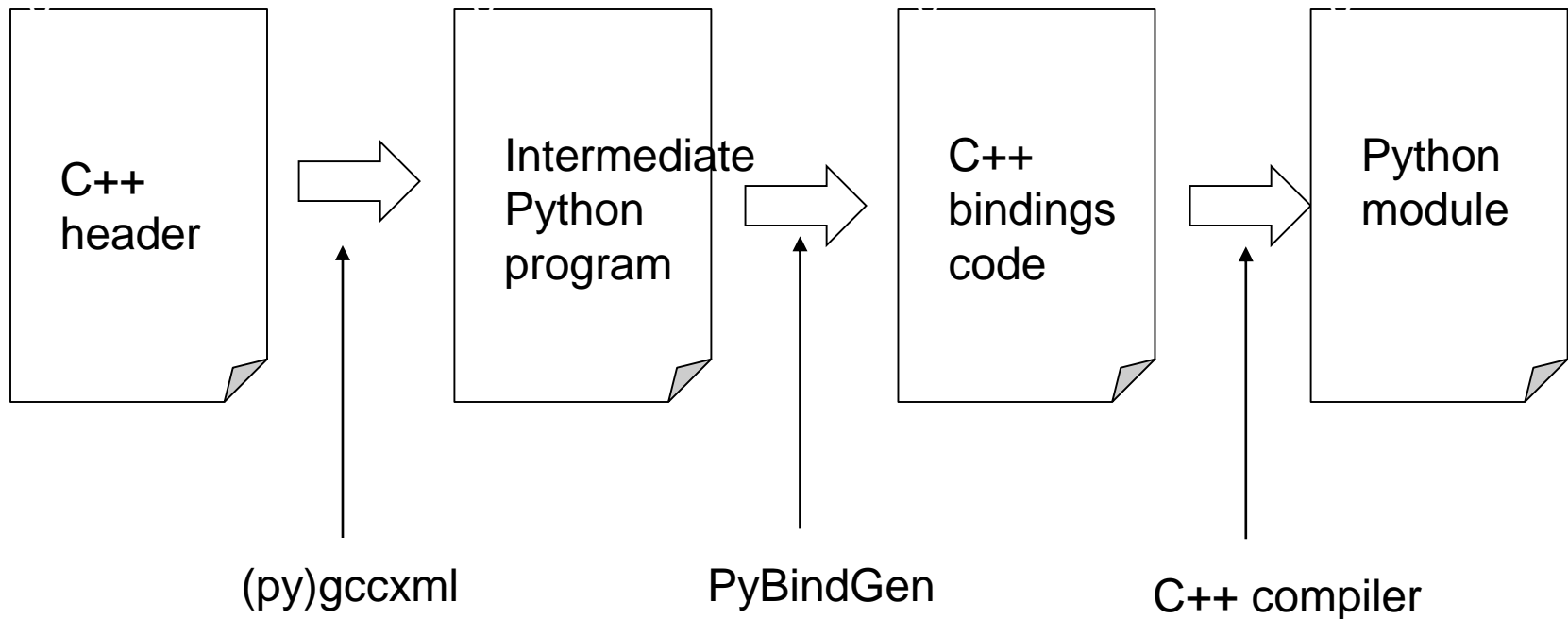


utilities



# Python bindings

- ns-3 uses a program called PyBindGen to generate Python bindings for all libraries



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# Why use ns-3?

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- You want to study *network performance* or *protocol operation* in a *controllable* or *scalable* environment
- You are comfortable writing C++ or Python code, and combining ns-3 with other code
- You like the idea of working on an active open source project
- ns-3 has the models you are looking for
  - or you can provide/integrate what is lacking

# What have people done with ns-3?

- ~750 publications to date
- search of 'ns-3 simulator' on IEEE and ACM digital libraries

1614 IEEE TRANSACTIONS ON NETWORKING, VOL. 20, NO. 4, DECEMBER 2012

## FSR: Formal Analysis and Implementation Toolkit for Safe Interdomain Routing

Andao Wang, Limin Jia, Member, IEEE, Wenchao Zhou, Yong Ruan, Boon Thau Loo, Member, IEEE, Senior Member, IEEE, Vivek Nigam, Andre Secevor, and Carolyn Talcott

**Abstract**—Interdomain routing stitches the disparate parts of the Internet together, enabling protocol stability, a critical issue to both researchers and practitioners. Yet, researchers create safety proofs and counterexamples by hand and build simulators and prototypes to explore protocol dynamics. Similarly, network operators analyze their router configurations manually or using homogenous tools. In this paper, we present a comprehensive toolkit for analyzing and implementing routing policies. Our high-level guidelines specify router configurations. Our *Formally Safe Routing (FSR)* toolset performs all of these tasks from the same algorithmic representation of routing policy. We show that routing algebra has a natural translation to both integer constraint routing safety analysis with SAT solvers and declarative programs (to generate distributed implementations). Our extensive experiment with realistic topologies and policies shows how FSR can detect problems in an autonomous system's (AS) BGP configuration, prove sufficient conditions for Border Gateway Protocol (BGP) safety, and empirically evaluate convergence time.

**Index Terms**—Communication technology, declarative networking, formal analysis, routing algebra.

### 1. INTRODUCTION

THE INTERNET'S global routing system does not necessarily converge, depending on how the Border Gateway Protocol (BGP) policies of individual networks are configured. Since protocol oscillations cause serious performance disruptions and route overhead, researchers devote significant attention to BGP stability (or "safety"). Almost a dozen models of BGP [1]–[15], [36] allow researchers to explore how local policies affect BGP stability and identify policy configurations that, if universally adopted by ISPs, ensure global convergence.

Given policy configurations as input, FSR produces an analysis of safety properties and a distributed protocol implementation, as shown in Fig. 1. FSR has three main underlying technologies:

- **Policy configuration as algebra:** Our extensions to routing algebra [13], [36] allow researchers and network operators to express policy configurations in an abstract algebraic form. These configurations can be anything from high-level policy guidelines (e.g. proposed counterfactuals that a researcher wants to study) or a completely specified policy instance (e.g. an ASBGP configuration or a multi-autonomous system's (AS) network that an operator wants to analyze). Route configurations can be automatically translated into the algebraic representation, easing the adoption of FSR.
- **Safety analysis:** To automatically analyze the policy configuration, FSR reduces the convergence proof to a

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**Abstract** We present an efficient message delivery framework, called *MedDa*, which enables communication in an Internet connecting heterogeneous networks that is prone to disruptions in connectivity. *MedDa* is complementary to the IETF's Bundle Architecture: besides its ability to store messages for unavailable destinations, *MedDa* can bridge the connectivity gap between infrastructure-based and multi-hop infrastructure-less networks. It benefits from network heterogeneity (e.g., nodes supporting more than one network and nodes having diverse resources) to improve message delivery. For example, in IEEE 802.11 networks, participating nodes can use both infrastructure and ad-hoc modes to deliver data to otherwise unavailable destinations. It also employs opportunistic routing to support nodes with episodic connectivity. One of *MedDa*'s key features is that any *MedDa* node can relay data to any destination and can act as a gateway to make two networks inter-operate or to connect to the backbone network. The network is able to store data destined to temporarily unavailable nodes in the time of their expiry. This time period depends upon current storage availability as well as quality-of-service needs (e.g., delivery delay bounds) imposed by the application. We showcase

**Keywords** Disruption tolerance · Episodic connectivity · Heterogeneous networks · Node relaying · Store-carry-and-forward · DTN routing

### 1. Introduction

It is envisioned that the Internet of the future will be highly heterogeneous not only due to the wide variety of end devices it interconnects, but also in terms of the underlying network it comprises. Figure 1 illustrates networks that range from wired and wireless backbones (e.g., commodity wireless mesh networks) to wireless infrastructure-based and ad-hoc networks (e.g., MANETs). On the other hand, current and emerging application, such as emergency response, environmental monitoring, smart environments (e.g., smart offices, homes, museums, etc.), and vehicular networks, among others, imply frequent and arbitrarily long-lived disruptions in connectivity. The resulting disruption- or delay-tolerant networks (DTNs) will likely become an important component of future internetworks.

Seamless interoperability among heterogeneous networks is a challenging problem as these networks may have very different characteristics. Node diversity may also

Manuscript received May 23, 2011; accepted January 21, 2012; approved by IEEE/ACM Transactions on Networking Editor Z. Li, March 28, 2012. Date of current version December 11, 2012. This work was supported in part by the NSF under Grant CCF-0820028, CNS-0835648, CNS-0941132, CNS-1264970, CNS-0923287, IS-1121270, and TC-0805667; the ACORN under Grant FA9550-09-1-0041; the ONR under Grant N00014-09-1-0770 and N00014-11-1-0515; a gift from Intel; and the Air Force Office of Scientific Research.

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Wireless New (2011) 17:1735-1744  
DOI 10.1007/s11264-011-0377-0

## Message delivery in heterogeneous networks prone to episodic connectivity

Rao Naved Bin Rais · Thierry Turletti · Kati Orosz

Published on 17 August 2011  
© Springer Science+Business Media, LLC 2011

**Abstract** We present an efficient message delivery framework, called *MedDa*, which enables communication in an Internet connecting heterogeneous networks that is prone to disruptions in connectivity. *MedDa* is complementary to the IETF's Bundle Architecture: besides its ability to store messages for unavailable destinations, *MedDa* can bridge the connectivity gap between infrastructure-based and multi-hop infrastructure-less networks. It benefits from network heterogeneity (e.g., nodes supporting more than one network and nodes having diverse resources) to improve message delivery. For example, in IEEE 802.11 networks, participating nodes can use both infrastructure and ad-hoc modes to deliver data to otherwise unavailable destinations. It also employs opportunistic routing to support nodes with episodic connectivity. One of *MedDa*'s key features is that any *MedDa* node can relay data to any destination and can act as a gateway to make two networks inter-operate or to connect to the backbone network. The network is able to store data destined to temporarily unavailable nodes in the time of their expiry. This time period depends upon current storage availability as well as quality-of-service needs (e.g., delivery delay bounds) imposed by the application. We showcase

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## Augmenting Data Center Networks with Multi-Gigabit Wireless Links

Daniel Halperin<sup>1</sup>, Srikanth Kandula<sup>1</sup>, Jitendra Padhye<sup>1</sup>, Paramvir Bahl<sup>1</sup>, and David Wetherall<sup>1</sup>  
Microsoft Research<sup>1</sup> and University of Washington

**Abstract**—The 60GHz wireless technology that is now emerging has the potential to provide dense and extremely fast connectivity at low cost. In this paper, we explore its use to relieve hotspots in over-subscribed data center (DC) networks. By experimenting with prototype equipment, we show that DC environments are well suited to a deployment of 60GHz links contrary to concerns about interference and link reliability. Using directional antennas, many wireless links can run concurrently at multi-Gbps rates on top-of-rack (ToR) switches. The wired DC network can be used to sidestep several common wireless problems. By analyzing production traces of DC traffic for four real applications, we show that adding a small amount of network capacity in the form of wireless *flyways* to the wired DC network can improve performance. However, to be of significant value, we find that one hop indirect routing is needed. Informed by our 60GHz experiments and DC traffic analysis, we present a design that uses DC traffic levels to select and add flyways to the wired DC network. Trace-driven evaluations show that network limited DC applications with predictable traffic workloads running on a 12 over-subscribed network can be sped up by 45% in 95% of the cases, with just one wireless device per ToR switch. With two devices, in 60% of the cases, the performance is identical to that of a non-over-subscribed network.

**Traditional wired DC networks** are un-structured and over-subscribed to keep costs down [15]. For example, a typical DC rack comprises 40 machines connected to a top-of-the-rack (ToR) switch with 1 Gbps links. The rack connects to an aggregation switch (to network with other racks) with 10 Gbps links. Thus, the link from the ToR to the aggregation switch can be over-subscribed with a ratio of 1/4. However, each over-subscribed link is a potential hotspot that hinders some DC application. Recent research tackles this problem by combining many more links and switches with variants of multihop routing so that the core of the network is no longer over-subscribed [1, 8, 9]. Of course, this benefit comes with large material cost and implementation complexity [15]. Some designs require many wire ties that cabling becomes a challenge [1], and most require ToR ToR links to be extra infrastructure.

In prior work [13], we argued instead for a more modest addition of links to relieve hotspots and boost application performance. The links, called *flyways*, add extra capacity to the core (by only a few ToR switches are used), a small number of flyways can significantly improve performance, without the cost of building a fully non-over-subscribed network.

The basic design of a DC network with 60GHz flyways is as follows. The base wired network is provisioned for the average case and can be over-subscribed. Each top-of-rack (ToR) switch is equipped with one or more 60GHz wireless devices, with electronic steerable directional antennas. A central controller monitors DC traffic patterns, and switches the beams of the wireless devices to set up flyways between ToR switches that provide added bandwidth as needed.

Other researchers have explored one of two other options and MEMS switches [7, 30] for steering flyways. We believe that 60GHz flyways are an attractive choice because wireless devices simplify because an important component of future internetworks.

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# Examples of recent publications

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- L. Salameh et al., "HACK: Hierarchical ACKs for Efficient Wireless Medium Utilization," *Proceedings of 2014 Usenix Annual Technical Conference* (best paper award winner), June 2014.
- A. Azgin et al., "Mobility Study for Named Data Networking in Wireless Access Networks," *Proceedings of IEEE ICC 2014*, June 2014.
- Wong S.-H and Gary Chan, "Topology Optimization for Wireless Mesh with Directional Antennas," *Proceedings of IEEE ICC 2014*, June 2014.
- C. Gouveia et al., "Development and implementation of Portuguese smart distribution system," *Electric Power Systems Research, Elsevier*, vol. 116, June 2014.
- M. Alharthi et al., "An Acumen/NS-3 integration for modeling networked Cyber-Physical Systems," *2014 Biennial Symposium on Communications (QSBC)*, June 2014.
- L. Ciarletta et al., "Simulation and platform tools to develop safe flock of UAVs: a CPS application-driven research," *2014 International Conference on Unmanned Aircraft Systems (ICUAS)*, May 2014.

# What have people done with ns-3?

- Educational use (from ns-3 wiki)

## Using ns-3 in Education

This page is a resource for learning about ns-3 as an educational tool for networking education.

## Papers

The [2011 Sigcomm Education workshop](#) had a paper regarding ns-3 use in the classroom:

- [An Open-source and Declarative Approach Towards Teaching Large-scale Networked Systems Programming](#)

## Courses using ns-3

The following courses have used ns-3 as courseware or to support projects

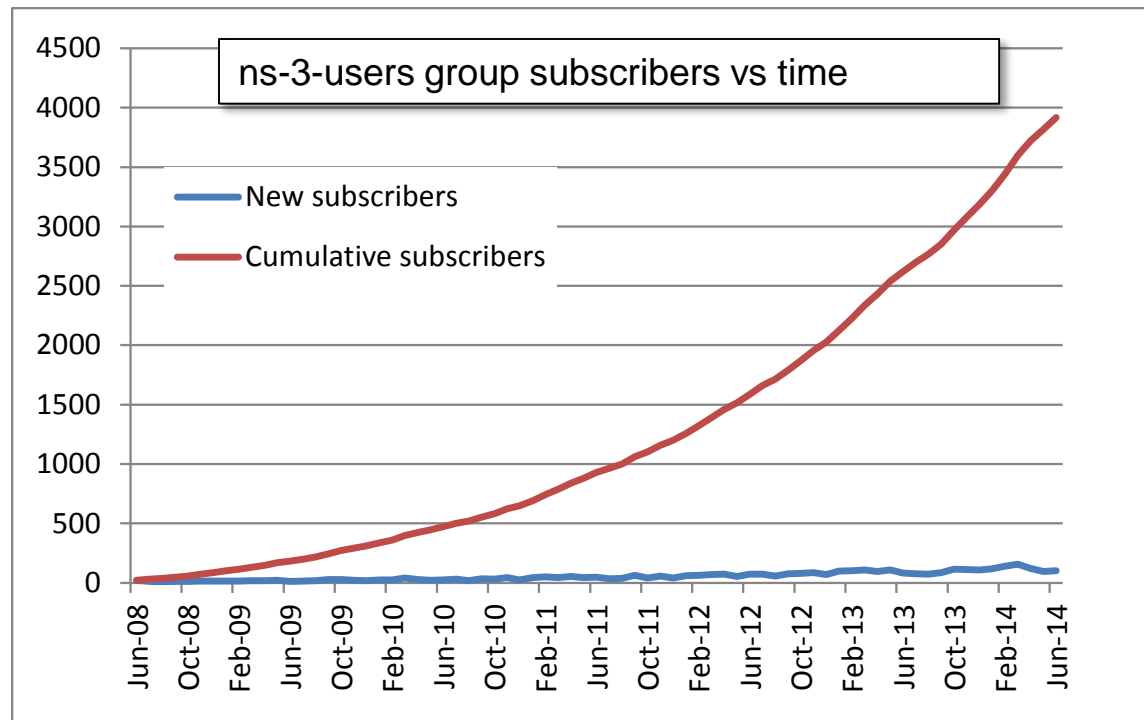
- [Georgia Tech. ECE 6110](#) Dr. George Riley, Spring 2013 (also Fall 2011, Fall 2010)
- [The University of Kansas EECS 780](#), [EECS 882](#), and [EECS 983](#) Dr. James Sterbenz, 2010 – 2012
- [UPenn CIS 553/TCOM 512](#) Dr. Boon Thau Loo, Fall 2010
- [Aalto University](#) Jose Costa-Requena and Markus Peuhkuri, Fall 2011
- [Indian Institute of Technology Bombay](#) Bhaskaran Raman, Autumn 2008
- [University of Rijeka](#)
  - [RM2-InfUniRi](#), Dr. Mario Radovan and [Vedran Miletić](#), Spring 2013, also Spring 2012
  - [RM-RiTeh](#), Dr. Mladen Tomić and [Vedran Miletić](#), Spring 2013

## Other resources

- Lalith Suresh's [Lab Assignments using ns-3](#) page.

# Statistics (July 2014)

- 3900 subscribers to ns-3-users
- 1440 subscribers to ns-developers
- ~ 15 maintainers
- ~ 150 authors/contributors





# Contributed code and associated projects

Overall ndnSIM documentation

Welcome to ndnSIM NS-3 based NDN simulator

We invite you to [join our mailing list](#) to see and participate in discussions about ndnSIM implementation and simulations in general ([mailing list archives](#)).

Contents:

- Introduction
  - More documentation
  - Support
  - A very short guide to the code
  - Logging
- Getting Started
  - Portability
  - Requirements
  - Downloading ndnSIM source
  - Compiling and running ndnSIM
- ndnSIM helpers
  - Stackbuilder
    - Routing
      - Manually routes
      - Global routing controller
      - Default routes
    - Content Store
    - Pending Interest Table
    - Forwarding strategy
  - Application
  - Content Store
    - Simple content stores
      - Least Recently Used (LRU) (default)
      - First-In-First-Out (FIFO)
      - Random
    - Content stores with entry lifetime tracking
      - Least Recently Used (LRU)
      - First-In-First-Out (FIFO)
      - Random
    - Content stores respecting freshness field of ContentObjects
      - Least Recently Used (LRU)
      - First-In-First-Out (FIFO)
      - Random

mptcp-ns3  
implement multipath TCP on ns-3

Project Home Downloads Wiki Issues Source

Summary Updates People

### Project description

The mptcp-ns3 project focuses on developing implementation of Multipath TCP on ns-3 for research purposes. The project implement the entire transport layer in ns-3.

Multipath TCP is an extension to TCP which aims to use multiple paths to handle a communication between two endpoints. MPTCP is the IETF working group to standardize Multipath TCP.

Please check the following URL for more information about multipath TCP: <http://datatracker.ietf.org/wg/mptcp/charter/>

### Current Status

The current implementation is really close to the MPTCP specification:

- MPTCP options: *MPC* (Multipath Capable), *ADD* and *REMOVE* address, *JOIN*, etc.
- Congestion Control: *Fully Coupled*, *Uncoupled TCPs*, *Linked Increases*, *RTT Compensator*.
- Packet Reordering: None, Eifel, DSACK and F-RT0 algorithms

### Getting Started

Follow the instructions in the wiki page <http://code.google.com/p/mptcp-ns3/wiki/Makelt> to successfully run simulations.

Decentralized Systems and Network Services Research Group - TM & SCC

### PhysSimWiFi for NS-3

Contact: Jens Mitrag, Stylianos Papanastasiou (CSS)  
Project: DSN, Chalmers University of Technology - Signals and Systems (CSS)  
Group:

### Overview

PhysSim-WiFi for [NS-3](#) is a detailed and accurate implementation of the OFDM-based IEEE 802.11 standard within the popular network simulator [NS-3](#). Compared to the default 802.11 PHY implementation of [NS-3](#), which abstracts packets by considering only an average signal strength per packet and the length of the packet, the PhysSim-WiFi implementation simulates the underlying signal processing steps of a transceiver down to the signal level, and introduces an increases accuracy for the decision whether a packet could be received correctly or not. At the same time, the new implementation allows to incorporate more sophisticated channel models. For instance, due to the modeling of packets on a signal level, channel models can emulate multi-path effects much more accurately and are able to reflect Doppler effects and their impact on the physical layer signal processing algorithm.

The PhysSim-WiFi implementation is a drop-in replacement of the default YansWiFPHY model, thus, it can be used with only minor modifications in the existing simulation code and the existing scenario setups.

For additional information and a changelog, please take a look at the

- [PhysSimWiFi Manual 1.1](#)

Full package download

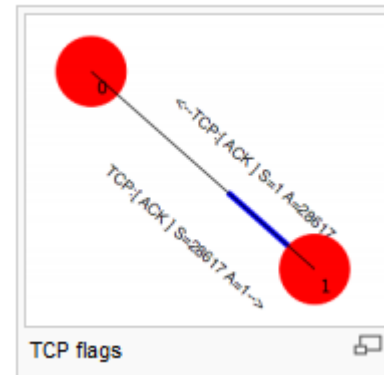
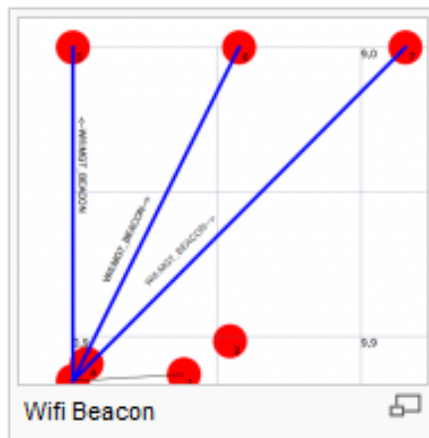
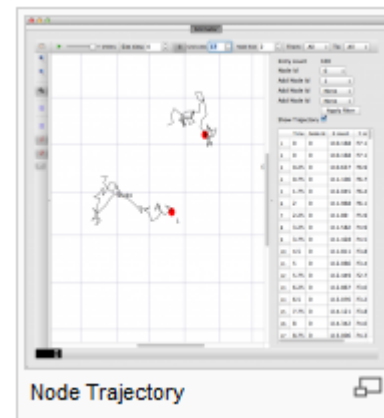
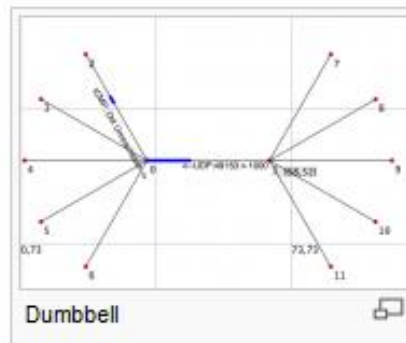
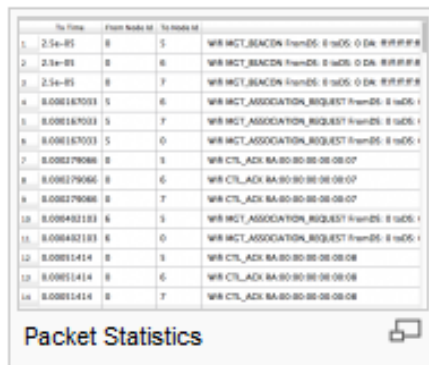
- [NS-3.9-PhysSimWiFi v1.1](#) (based on NS-3.9 and PhysSim-WiFi) - August 19, 2011
- [NS-3.9-PhysSimWiFi v1.0](#) (based on NS-3.9 and PhysSim-WiFi) - September 12, 2010

Patches for NS-3

- [PhysSimWiFi v1.1 for NS-3.9](#) - August 19, 2011
- [PhysSimWiFi v1.1 for NS-3.9-PhysSimWiFi v1.0](#) - August 19, 2011
- [PhysSimWiFi v1.0 for NS-3.9](#) - September 12, 2010

# NetAnim

- "NetAnim" by George Riley and John Abraham
  - see the 'ns3share' channel on YouTube

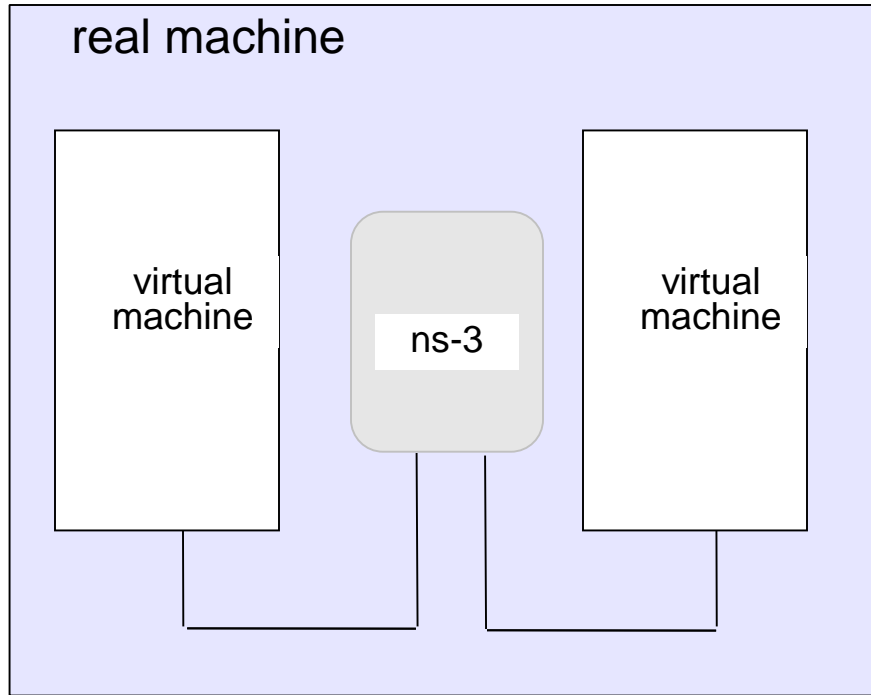


# Emulation support

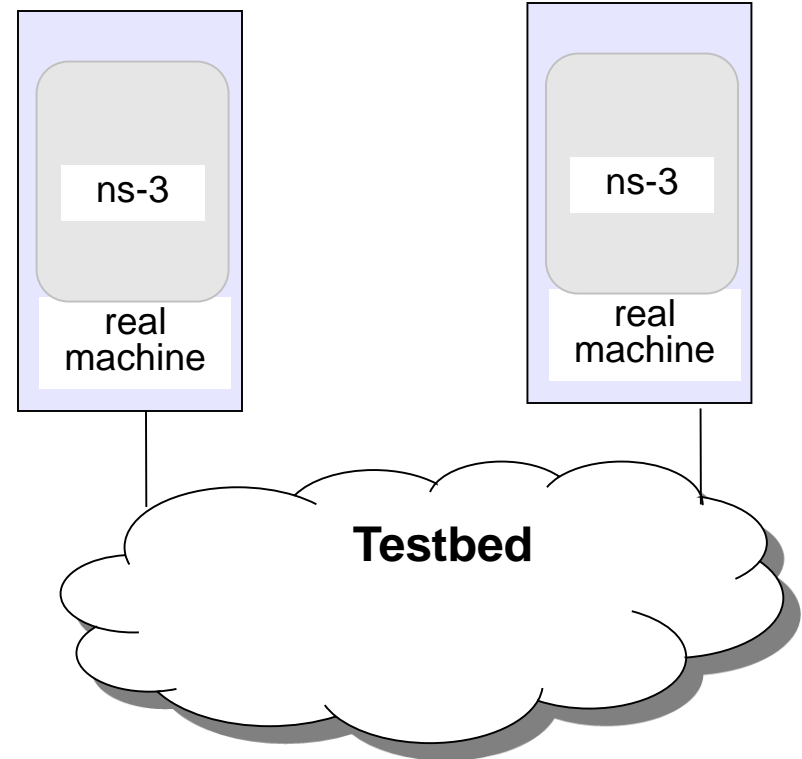
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- Support moving between simulation and testbeds or live systems
- A real-time scheduler, and support for two modes of emulation
- Linux is only operating system supported
- Must run simulator in real time
  - `GlobalValue::Bind ("SimulatorImplementationType", StringValue ("ns3::RealTimeSimulatorImpl"));`
- Must enable checksum calculations across models
  - `GlobalValue::Bind ("ChecksumEnabled", BooleanValue (true));`
- Must sometimes run as root

# ns-3 emulation modes



1) ns-3 interconnects real or virtual machines

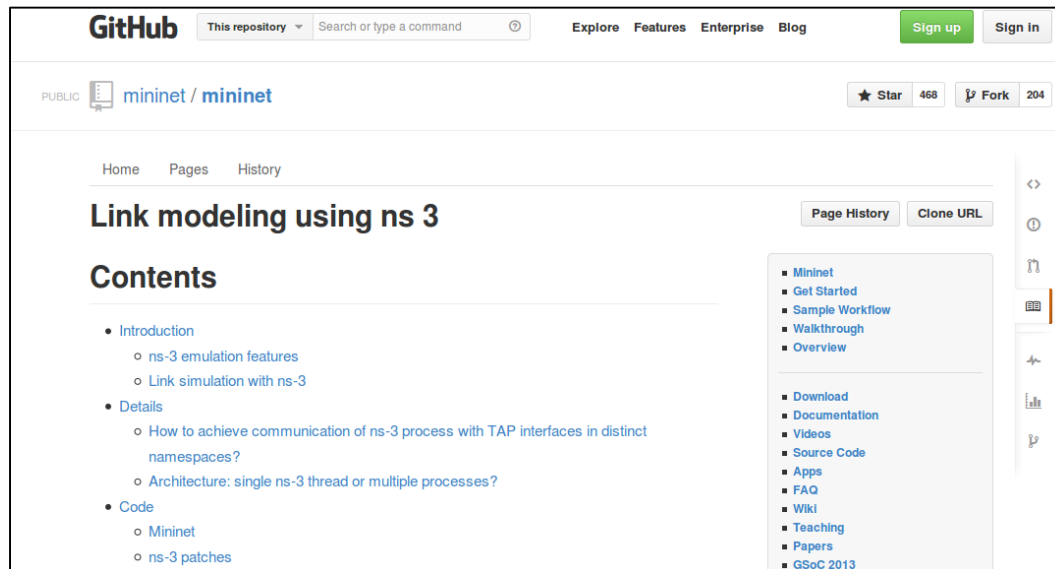


2) testbeds interconnect ns-3 stacks

Various hybrids of the above are possible

# Example use case: mininet

- Mininet is popular in the Software-Defined Networking (SDN) community
- Mininet uses "TapBridge" integration
- <https://github.com/mininet/mininet/wiki/Link-modeling-using-ns-3>



The screenshot shows the GitHub repository page for mininet/mininet. The page title is "Link modeling using ns 3". The "Contents" section lists the following items:

- Introduction
  - ns-3 emulation features
  - Link simulation with ns-3
- Details
  - How to achieve communication of ns-3 process with TAP interfaces in distinct namespaces?
  - Architecture: single ns-3 thread or multiple processes?
- Code
  - Mininet
  - ns-3 patches

The sidebar on the right contains a table of contents with the following items:

- Mininet
- Get Started
- Sample Workflow
- Walkthrough
- Overview
- Download
- Documentation
- Videos
- Source Code
- Apps
- FAQ
- Wiki
- Teaching
- Papers
- GSoC 2013

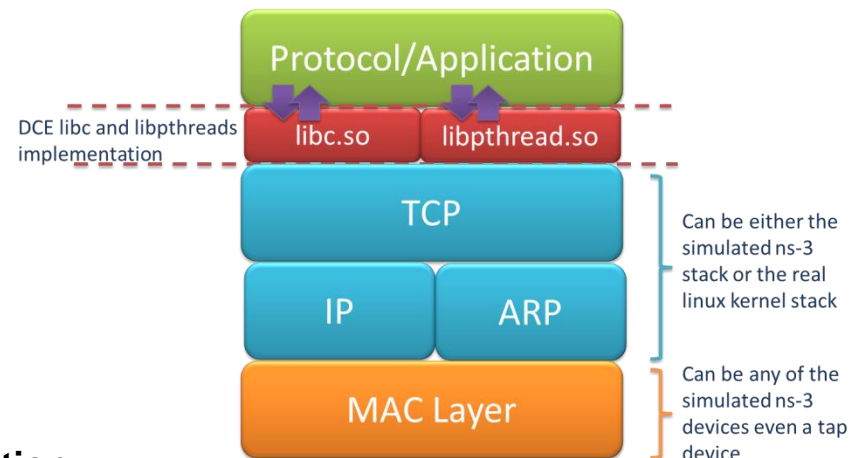
# Direct Code Execution

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- Lightweight virtualization of kernel and application processes, interconnected by simulated networks
- Benefits:
  - Implementation realism in controlled topologies or wireless environments
  - Model availability
  - Debugging a whole network within a single process
- Limitations:
  - Not as scalable as pure simulation
  - Tracing more limited
  - Configuration different

# Direct Code Execution implementation

- DCE/ns-3 framework requires the virtualization of a series of services
  - Multiple isolated instances of the same protocol on the same machine
- System calls are captured and treated by DCE
- Network stack protocols calls are captured and redirected
- To perform its work DCE re-implement the Linux program loader and parts of *libc* and *libpthread*



# DCE modes

- DCE modes in context of possible approaches

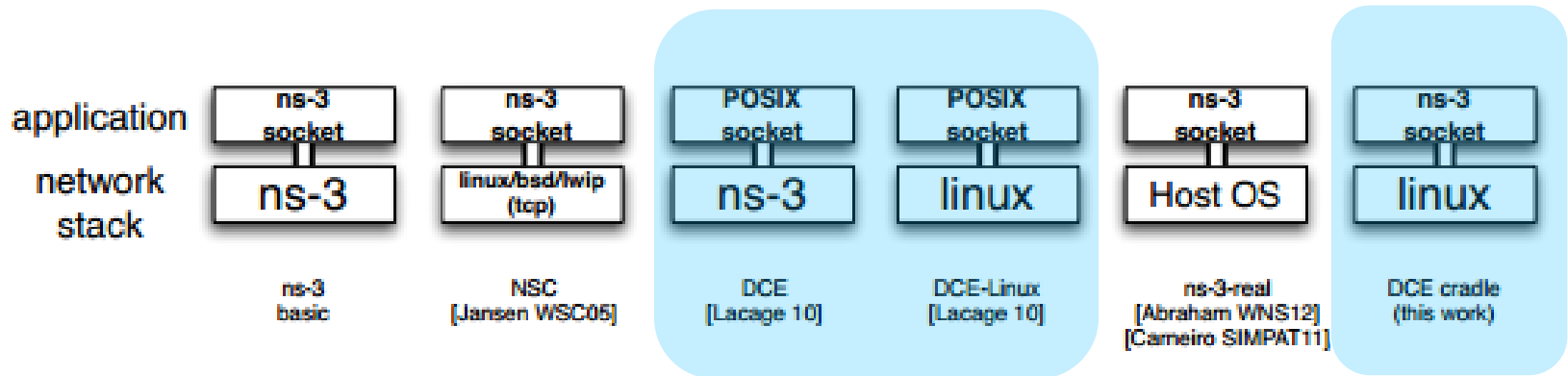


Figure 1: Current possible combinations of network stacks and applications.

Figure source: DCE Cradle: Simulate Network Protocols with Real Stacks for Better Realism, Tazaki et al, WNS3 2013.



# Agenda

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- ns-3 project overview
  - What is ns-3?
  - Why use ns-3?
  - Project organization
  - Relationship to ns-2
  - Future directions
- Getting started with ns-3

# *ns-3* project goals

---

Develop an extensible simulation environment for networking research

- 1) a tool aligned with the experimentation needs of modern networking research
- 2) a tool that elevates the technical rigor of network simulation practice
- 3) an open-source project that encourages community contribution, peer review, and long-term maintenance and validation of the software

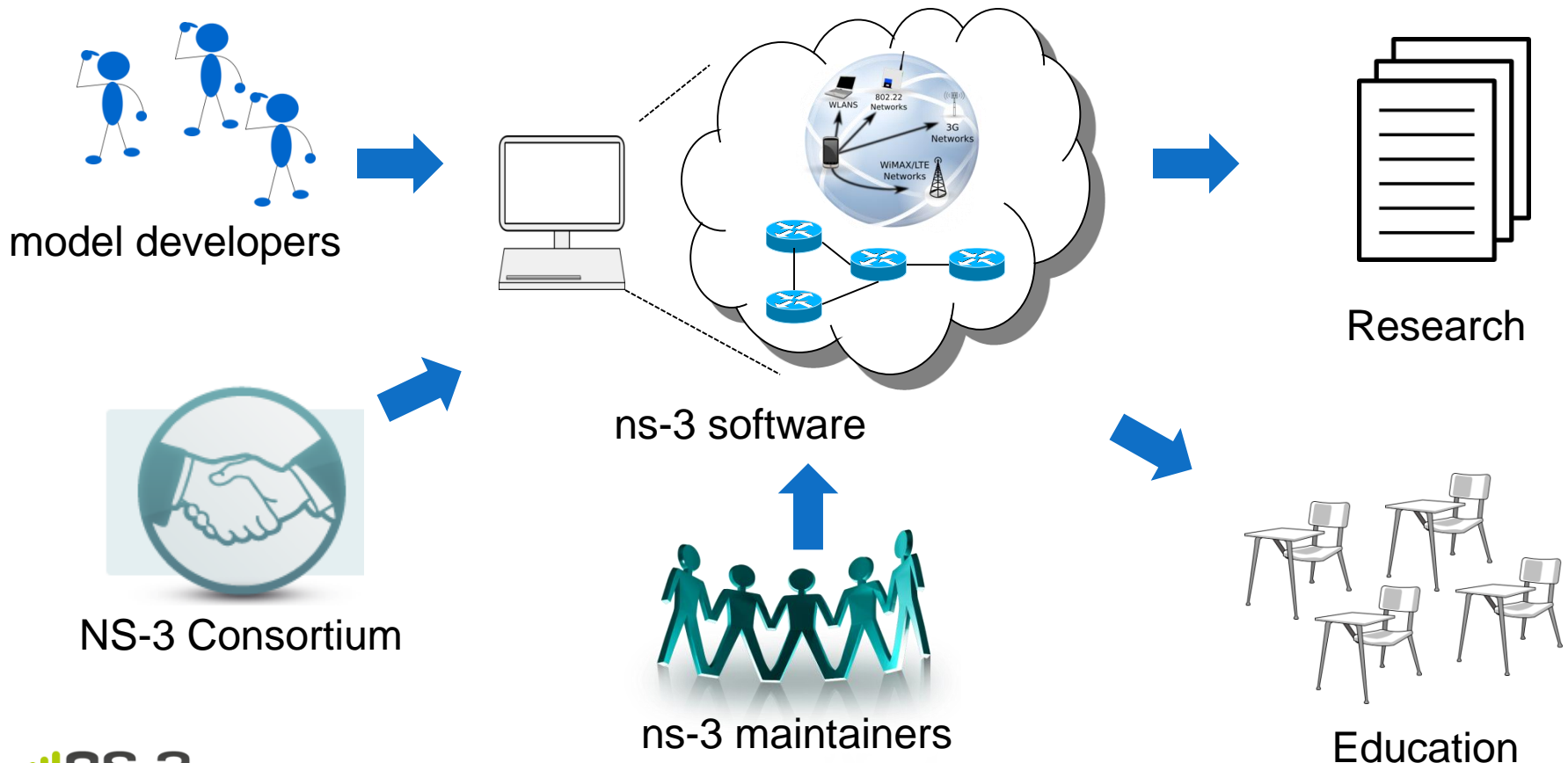
# How the project operates

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- Project provides three annual software releases
- Users interact on mailing lists and using Bugzilla bug tracker
- Code may be proposed for merge
  - Code reviews occur on a Google site
- Maintainers (one for each module) fix or delegate bugs, participate in reviews
- Project has been conducting annual workshop and developer meeting around SIMUTools through 2013
  - ns-3 Annual Meeting in Atlanta, May 2014
- Google Summer of Code (March-August) five of the past six summers

# ns-3: An Open Source Network Simulator

- ns-3 is a *discrete-event network simulator* targeted for *research and educational use*



# Goals of the NS-3 Consortium

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- The NS-3 Consortium is a collection of organizations cooperating to support and develop the ns-3 software.
- It operates in support of the open source project
  - by providing a point of contact between industrial members and ns-3 developers,
  - by sponsoring events in support of ns-3 such as users' days and workshops,
  - by guaranteeing maintenance support for ns-3's core, and
  - by supporting administrative activities necessary to conduct a large open source project.

# Acknowledgment of support



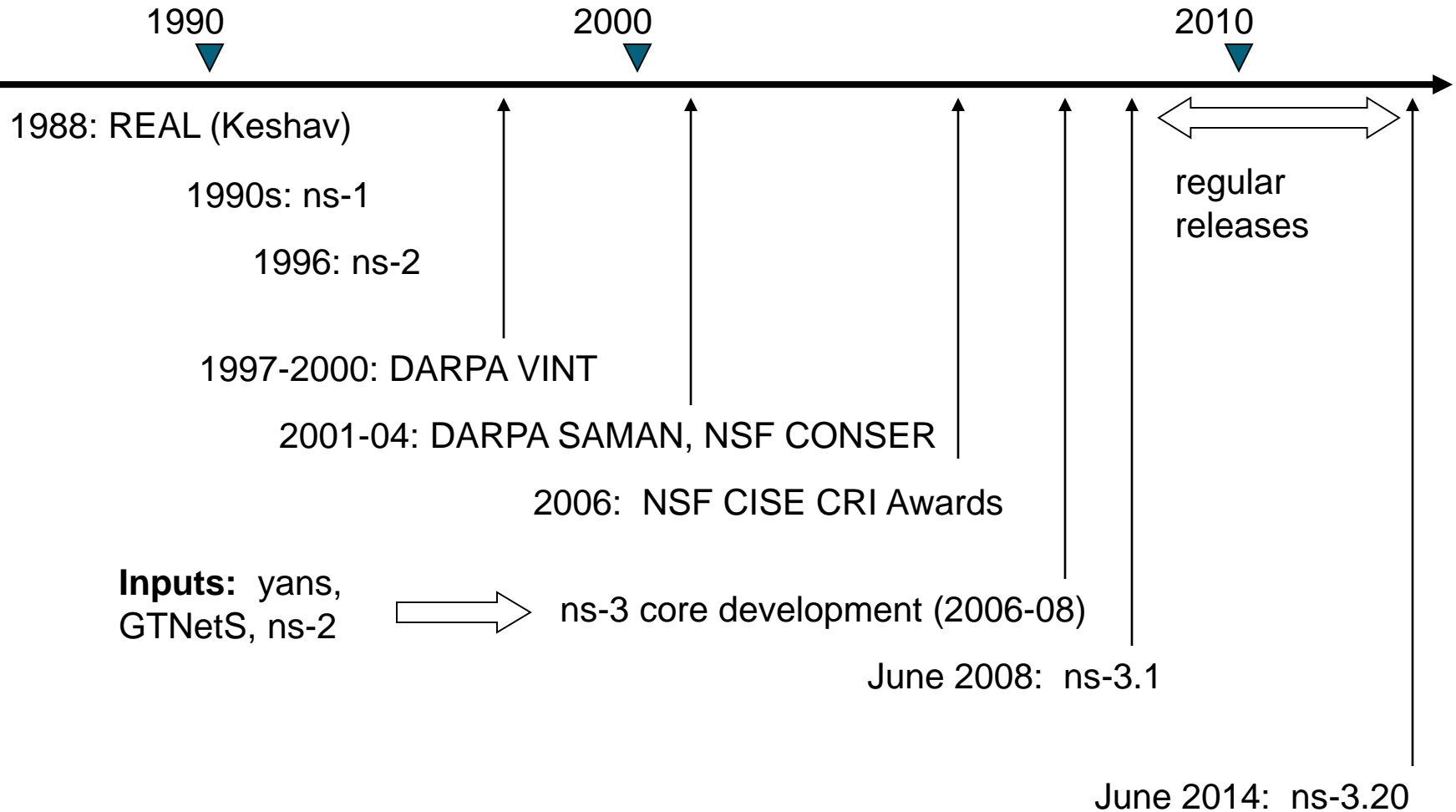
Information Sciences Institute

# Agenda

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# ns timeline





# Relationship to ns-2

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ns-3 is a new simulator, without backward compatibility

Similarities to ns-2:

- C++ software core
- GNU GPLv2 licensing
- ported ns-2 models: random variables, error models, OLSR, Calendar Queue scheduler

Differences:

- Python scripting (or C++ programs) replaces OTcl
- most of the core rewritten
- new animators, configuration tools, etc. are in work
- ns-2 is no longer actively maintained/supported

# Agenda

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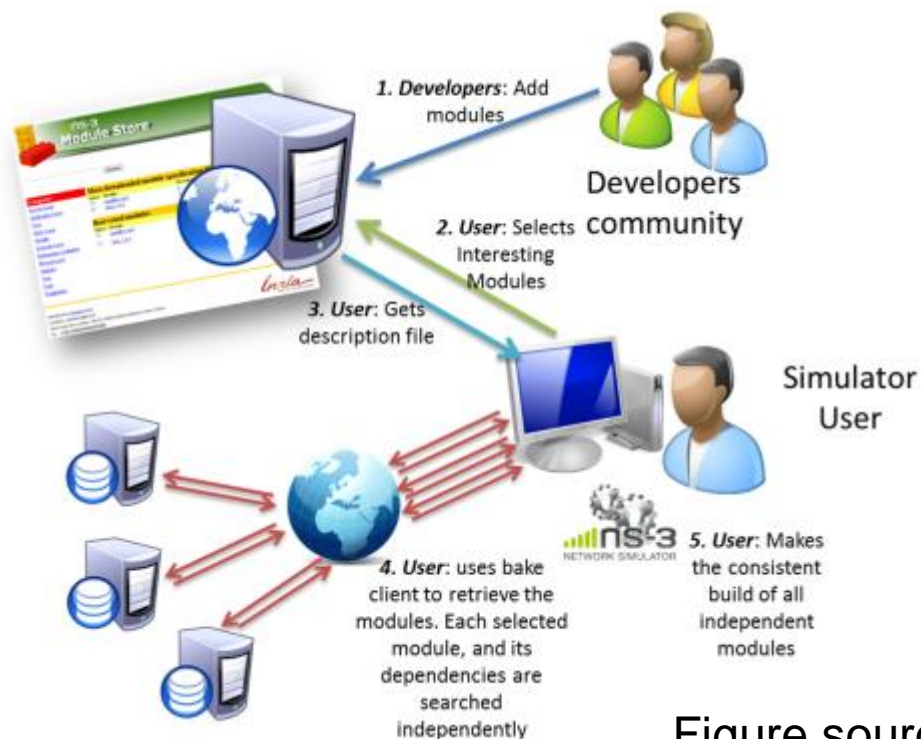
# Development Priorities

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- Software modularity and long-term maintenance
- Improved integration of direct code execution
- Improved integration with container-based and testbed-based experiment infrastructures
- Simulation-based experiment management
- Usability

# Modularity

- Open source project maintains a (more stable) core
- Models migrate to a more federated development process



"bake" tool (Lacage and Camara)

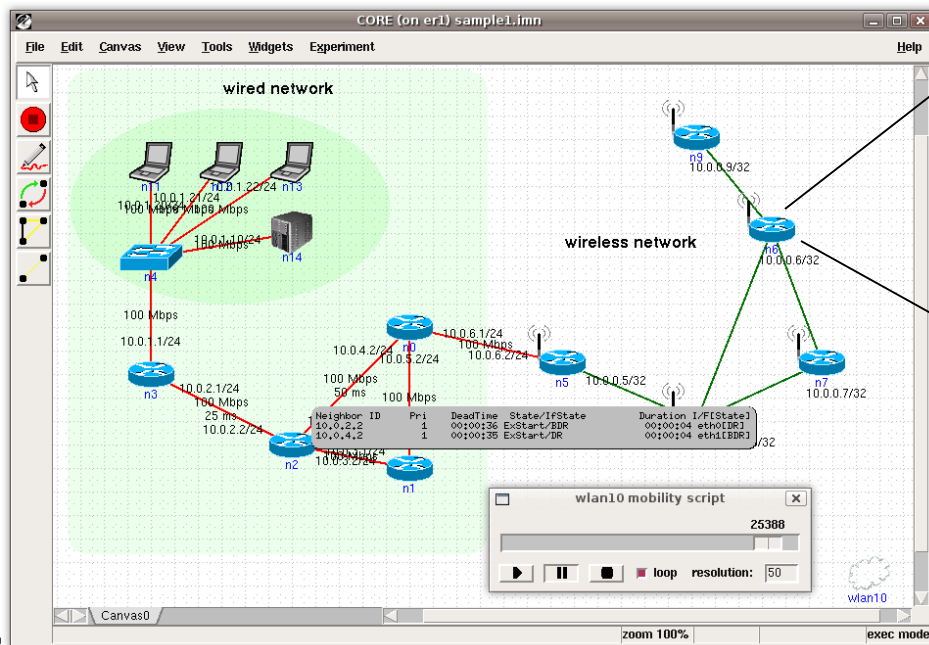
Components:

- build client
- "module store" server
- module metadata

Figure source: Daniel Camara

# Container-based Integration

- Common Open Research Emulator (CORE)
  - <http://pf.itd.nrl.navy.mil>
- Python-based framework using ns-3 Python bindings, distributed computing library, and ns-3 TapBridge framework



Object Attributes	Attribute Value
ns3::NodeListPriv	
NodeList	
0	
DeviceList	
0	
Address	00:00:00:00:00:01
EncapsulationMode	Llc
SendEnable	true
ReceiveEnable	true
DataRate	5000000bps
TxQueue	
1	
ApplicationList	
ns3::PacketSocketFactory	
ns3::Ipv4L4Demux	
ns3::Tcp	
ns3::Udp	
ns3::Ipv4	
ns3::ArpL3Protocol	
ns3::Ipv4L3Protocol	

Figure source:  
Jeff Ahrenholz

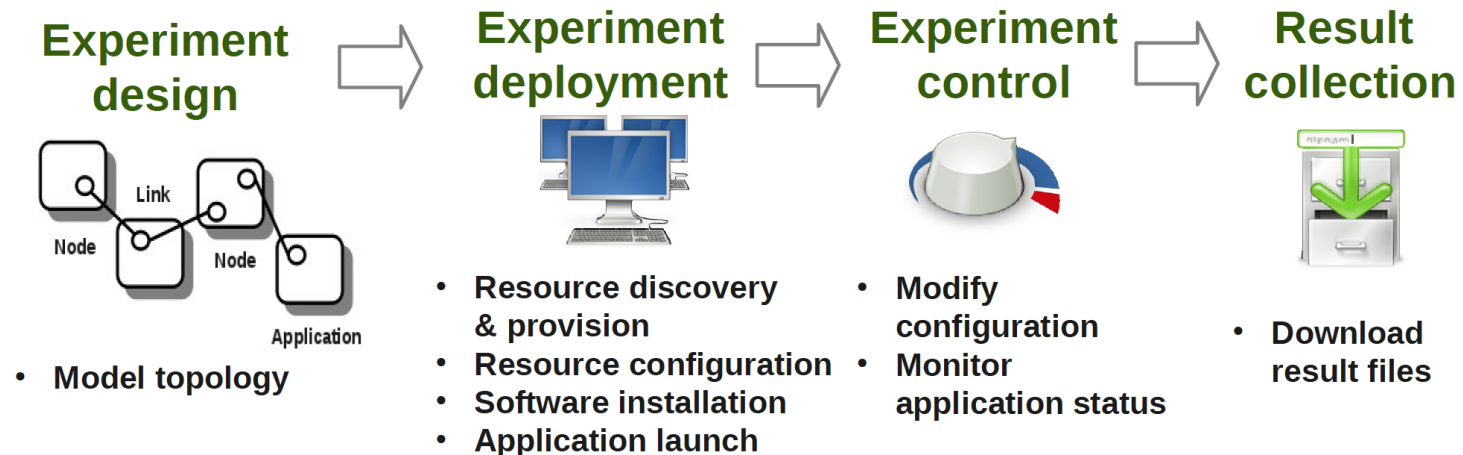
# General issues with hybrid environments

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- Ease of use
  - Configuration management and coherence
  - Information coordination (two sets of state)
    - e.g. IP/MAC address coordination
  - Output data exists in two domains
  - Debugging
- Error-free operation (avoidance of misuse)
  - Synchronization, information sharing, exception handling
    - Checkpoints for execution bring-up
    - Inoperative commands within an execution domain
    - Deal with run-time errors
  - Soft performance degradation (CPU) and time discontinuities

# Network Experiment Management Framework (NEPI)

- Network experiment management framework to automate experiment life-cycle
- Allows scenarios involving heterogeneous resources (ns-3, PlanetLab, netns, ...)
- Wiki: <http://nepi.inria.fr>



*Figure source: Alina Quereilhac, INRIA*

# SAFE: Simulation Automation Framework

- Data collection, transient analysis, management of independent replications, graphical configuration and visualization
- In ns-2 realm, similar to projects like ANSWER, ns2measure, and Akaroa2

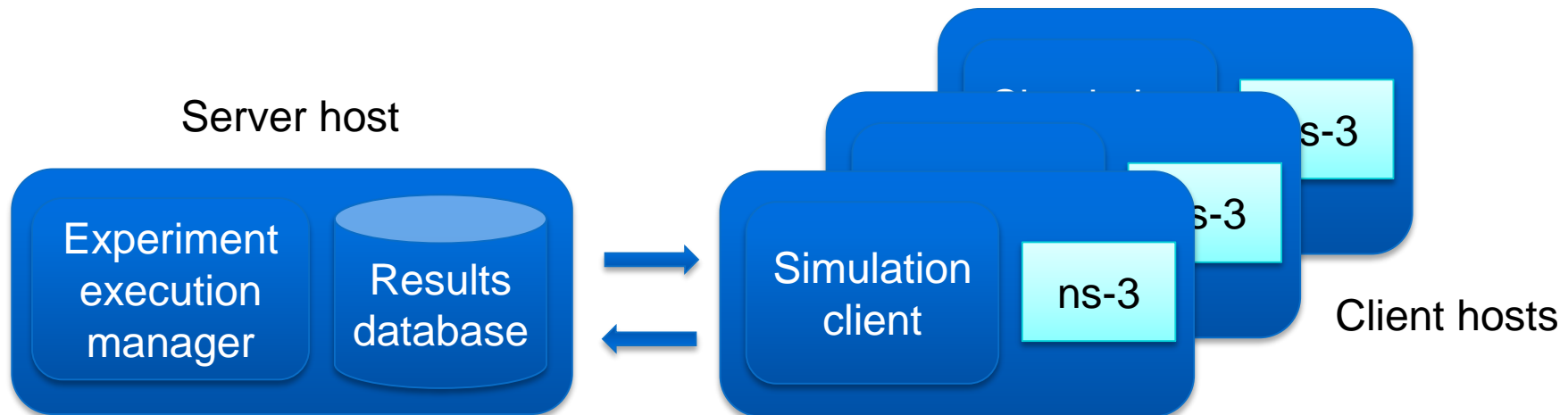
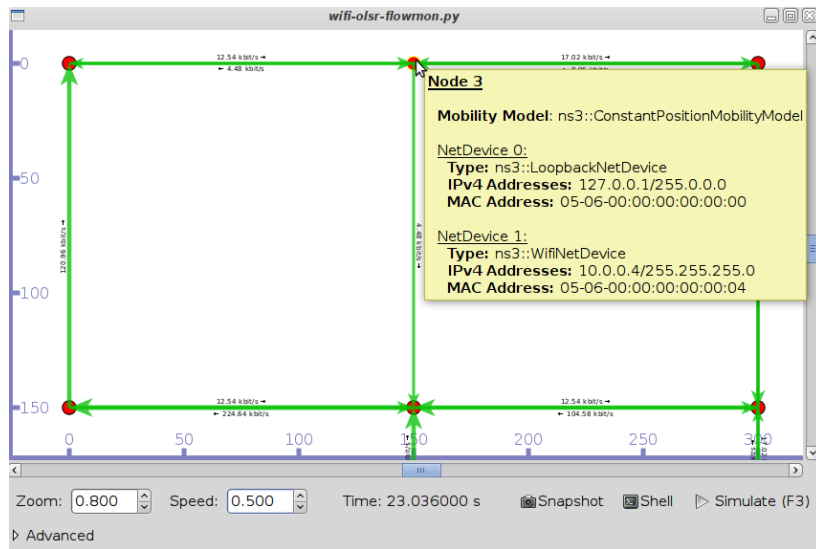


Figure source: Felipe Perrone

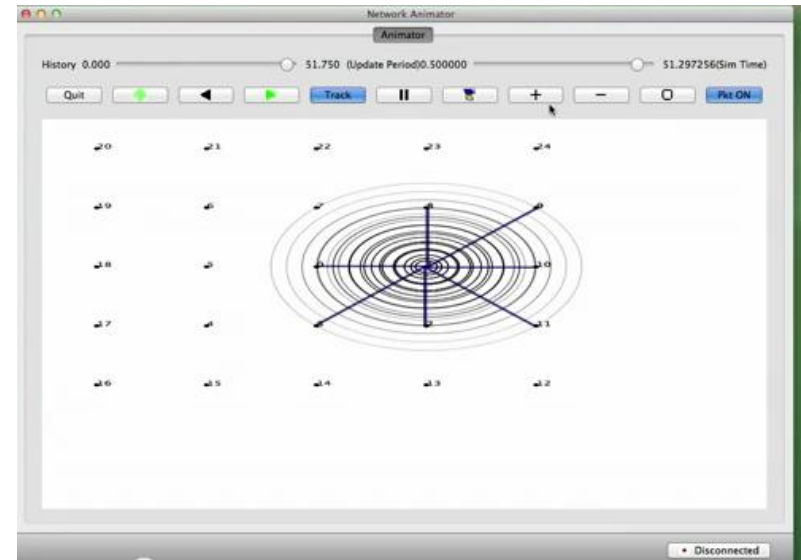


# Usability

- Animation and visualization



PyVis (Carneiro)



NetAnim (Riley and Abraham)

- Linkage to external tools (topology, mobility, statistics)
- Improved helper APIs

# Agenda

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# Getting started with ns-3

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- Finding what you need
- Contributing to the project

# Resources

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Web site:

<http://www.nsnam.org>

Mailing lists:

<https://groups.google.com/forum/#!forum/ns-3-users>

<http://mailman.isi.edu/mailman/listinfo/ns-developers>

Wiki:

<http://www.nsnam.org/wiki/>

Tutorial:

<http://www.nsnam.org/docs/tutorial/tutorial.html>

IRC: #ns-3 at freenode.net

# Suggested steps

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- Work through the ns-3 tutorial
- Browse the source code and other project documentation
  - manual, model library, Doxygen, wiki
  - ns-3 Consortium tutorials (March 2013)
    - <http://www.nsnam.org/consortium/activities/annual-meeting-march-2013/>
- Ask on ns-3-users mailing list if you still have questions
  - We try to answer most questions

# APIs

- Most of the ns-3 API is documented with Doxygen
  - <http://www.stack.nl/~dimitri/doxygen/>

NS-3

- ns-3 Documentation
- NS-3 Modules
- NS-3 Class List
- NS-3 Class Hierarchy
- Class Members
- NS-3 Graphical Class Hierarchy
- NS-3 Namespace List
- Namespace Members
- NS-3 Related Pages

Main Page Modules Namespaces Classes Related Pages

Class List Class Hierarchy Class Members

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ns3::InetSocketAddress

---

### ns3::InetSocketAddress Class Reference

[Address]

an Inet address class [More...](#)

```
#include <inet-socket-address.h>
```

Collaboration diagram for ns3::InetSocketAddress:

```
graph TD; ns3::InetSocketAddress -- m_ipv4 --> ns3::Ipv4Address;
```

[Legend]

[List of all members.](#)

Public Member Functions

# Reading existing code

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- Much insight can be gained from reading ns-3 examples and tests, and running them yourselves
- Many core features of ns-3 are only demonstrated in the core test suite (src/core/test)
- Stepping through code with a debugger is informative
  - callbacks and templates make it more challenging than usual

# FAQs

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- Does ns-3 have a Windows version?
  - Yes, for Visual Studio 2012
  - [http://www.nsnam.org/wiki/Ns-3\\_on\\_Visual\\_Studio\\_2012](http://www.nsnam.org/wiki/Ns-3_on_Visual_Studio_2012)
- Does ns-3 support Eclipse or other IDEs?
  - Instructions have been contributed by users
  - [http://www.nsnam.org/wiki/HOWTO\\_configure\\_Eclipse\\_with\\_ns-3](http://www.nsnam.org/wiki/HOWTO_configure_Eclipse_with_ns-3)
- Is ns-3 provided in Linux or OS X package systems (e.g. Debian packages)?
  - No



# Contributing

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- Any amount of help is appreciated!
  - Reporting stale documentation to [webmaster@nsnam.org](mailto:webmaster@nsnam.org)
  - Contributing small patches
  - Writing new documentation
  - Reporting bugs
  - Fixing bugs
  - Reviewing code of others
  - Contributing new code
  - Becoming a maintainer

# New project ideas

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- Visit the wiki under "Project Ideas" tab
  - [http://www.nsnam.org/wiki/Project\\_Ideas](http://www.nsnam.org/wiki/Project_Ideas)
- Students, consider to apply for Google Summer of Code 2015
  - A 10-week summer job that mentors a student project on ns-3
  - Students apply in March 2015 timeframe
  - <http://www.nsnam.org/wiki/GSOC2014Projects>

# Questions?

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